

Introduction to High-Performance Parallel Distributed Computing using Chapel, UPC++, and Coarray Fortran

ECP/NERSC/OLCF 2023 Tutorial 30-minute Intro session go.lbl.gov/cuf23



Hewlett Packard Enterprise



CAK RIDGE National Laboratory



National Energy Research Scientific Computing Center



EXASCALE COMPUTING PROJECT

Introduction to High-Performance Parallel Distributed Computing using Chapel, UPC++ and Coarray Fortran



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Acknowledgements

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Schedule for Chapel, UPC++ and Fortran Tutorial

Wed July 26, noon - 3:15pm (all times US Eastern)

- noon 1:30: Tutorial Overview
 - including a 20-minute intro to each programming model
- 1:30 1:45: Coffee Break
- 1:45 3:15: Parallel programming in Chapel

Thu July 27, noon - 3:15pm

- noon 1:30: Parallel programming with UPC++
- 1:30 1:45: Coffee Break
- 1:45 3:15: Parallel programming with Fortran Coarrays







Audience questions

Slack is preferred: go.lbl.gov/cuf23-slack

alternatively use Zoom chat

Motivation

- You have ...
 - A lot of data to process and analyze
 - A big simulation to run
 - Or both of the above
- Resources are available
 - Your laptop has multiple cores that can process in parallel
 - Your lab/institution has a cluster
 - Or your lab/institution has a supercomputer
- Writing a parallel program enables you to analyze data and/or perform simulations significantly faster.







When poll is active, respond at pollev.com/michellestrout402
 Text MICHELLESTROUT402 to 22333 once to join

Which programming language(s) do you use the most? (you can respond to this question 3 times)





PGAS Programming Models

• PGAS: Partitioned Global Address space

CHAPEL

- Chapel, UPC++, and Fortran with coarrays are PGAS programming models
- A programming model provides an interface and code patterns to a programmer along with a concept of how code will execute at runtime.

PGAS Programming Models

- Can access variables in global address space from each node
- Implemented with puts and gets (RMA: remote memory access)
- Can partition/organize data and computation to reduce RMA





This tutorial: Chapel, UPC++, Fortran with coarrays

- Shared example shown in all three: **2D heat diffusion**
- Then other examples per programming model
 - Chapel: k-mer counting, image analysis, processing files in parallel
 - UPC++: 1-d Jacobi solver, distributed hash table
 - Fortran: 2-d heat equation, hello world variants
- Hands On
 - Providing a cloud instance, Perlmutter, and Frontier instructions for obtaining a tarball containing all example programs: <u>go.lbl.gov/cuf23</u>
 - You are encouraged to compile, run, and experiment with the examples throughout
- Q&A Protocol
 - Model experts are available to answer questions in Slack: **<u>go.lbl.gov/cuf23-slack</u>**
 - You should have received an email invite, or can follow the link above





Production Applications using these Programming Models



CHAMPS: 3D Unstructured CFD

(~100K lines of Chapel) Éric Laurendeau, Simon Bourgault-Côté, Matthieu Parenteau, et al. École Polytechnique Montréal



ICAR: Intermediate Complexity Atmospheric Research model written in Coarray Fortran

https://github.com/NCAR/icar

MetaHipMer, a genome assembler written in UPC++



Hands On: Compiling and Running Hello Worlds

- Instructions on how to compile and run a **hello world** for all three programming models.
- Hands-on examples and instructions: go.lbl.gov/cuf23
 - Options include:
 - NERSC Perlmutter, OLCF Frontier, AWS Cloud, Docker, ...
 - Pause here for attendees to setup their programming environment





Do you have any parallel programming experience? If so, what tools have you used?



Start the presentation to see live content. For screen share software, share the entire screen. Get help at pollev.com/app

Shared Problem: 2D Heat Diffusion



- Specifically a 2D heat diffusion problem
 - 2D diffusion equation is above. Mathematical details: <u>wikipedia.org/wiki/Heat_equation</u>
 - Discretization solving for the unknown at time step n+1 and spatial coordinate i,j
- Steps in sample codes
 - \circ Set some initial conditions for u⁰
 - Estimate u over time and space as shown below
 - Show how to parallelize these computations

$$u_{i,j}^{n+1} = u_{i,j}^{n} + \frac{\nu \Delta t}{\Delta x^{2}} (u_{i+1,j}^{n} - 2u_{i,j}^{n} + u_{i-1,j}^{n}) \qquad \frac{\text{Simplified form}}{\text{assume } \Delta \mathbf{x} = \Delta \mathbf{y}, \text{ and let } \alpha = \nu \Delta t / \Delta \mathbf{x}^{2}} \\ + \frac{\nu \Delta t}{\Delta y^{2}} (u_{i,j+1}^{n} - 2u_{i,j}^{n} + u_{i,j-1}^{n}) \qquad u_{i,j}^{n+1} = u_{i,j}^{n} + \alpha \begin{pmatrix} u_{i+1,j}^{n} + u_{i-1,j}^{n} \\ -4u_{i,j}^{n} + u_{i,j+1}^{n} + u_{i,j-1}^{n} \end{pmatrix}$$

Three questions about how you program

- Have you used a cluster or supercomputer before? If so, what were their characteristics (number of nodes, threads per node, etc)?
- Where do you go when you have programming questions? A colleague, stack overflow, google search, documentation, ...
- For your code, what computations/libraries are most important for your work?

NOTE: The pollEV survey starts on the next slide, but it won't show the above questions. This slide is to show you what those questions will be.









Three questions about how you program



0 done

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Start the presentation to see live content. For screen share software, share the entire screen. Get help at pollev.com/app

What do you want to learn about Chapel, UPC++, or Coarray Fortran today?

Тор



Start the presentation to see live content. For screen share software, share the entire screen. Get help at pollev.com/app

Schedule for Chapel, UPC++ and Fortran Tutorial

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 - Chapel Intro
 - Fortran with co-arrays Intro
 - UPC++ Intro
- 1:30 1:45: Coffee Break
- 1:45 3:15: Parallel programming in Chapel

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go.lbl.gov/cuf23

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INTRODUCTION TO CHAPEL PARALLEL PROGRAMMING LANGUAGE

Michelle Strout and Jeremiah Corrado CUF23: Sponsored by OLCF, NERSC, and ECP July 26-27, 2023

INTRODUCTION TO CHAPEL

- What Chapel is and how programmers are using Chapel in their applications
- Chapel execution model with a parallel and distributed "Hello World"
- 2D Heat Diffusion example: variants and how to compile and run them
- Learning objectives for today's 90-minute Chapel tutorial

CHAPEL PROGRAMMING LANGUAGE

Chapel is a general-purpose programming language that provides ease of parallel programming, high performance, and portability.

And is being used in applications in various ways:

refactoring existing codes,

developing new codes,

serving high performance to Python codes (Chapel server with Python client), and providing distributed and shared memory parallelism for existing codes.

APPLICATIONS OF CHAPEL: LINKS TO USERS' TALKS (SLIDES + VIDEO)



CHAMPS: 3D Unstructured CFD
CHIUW 2021 CHIUW 2022



Arkouda: Interactive Data Science at Massive Scale CHIUW 2020 CHIUW 2023



CHIUW 2021 CHIUW 2023



ChplUltra: Simulating Ultralight Dark MatterCHIUW 2020CHIUW 2022





Lattice-Symmetries: a Quantum Many-Body Toolbox Desk dot chpl: Utilities for Environmental Eng.





Chapel-based Hydrological Model Calibration



CHIUW 2022



CrayAl HyperParameter Optimization (HPO)



RapidQ: Mapping Coral Biodiversity

CHIUW 2023



CHGL: Chapel Hypergraph Library CHIÚW 2020



ChapQG: Layered Quasigeostrophic CFD



Your Application Here?



HIGHLIGHTS OF CHAPEL USAGE

CHAMPS: Computational Fluid Dynamics framework for airplane simulation

- Professor Eric Laurendeau's team at Polytechnique Montreal
- Performance: achieves competitive results w.r.t. established, world-class frameworks from Stanford, MIT, etc.
- Programmability: "We ask students at the master's degree to do stuff that would take 2 years and they do it in 3 months."

Arkouda: data analytics framework (<u>https://github.com/Bears-R-Us/arkouda</u>)

- Mike Merrill, Bill Reus, et al., US DOD
- Python front end client, Chapel server that processes dozens of terabytes in seconds
- April 2023: 1200 GiB/s for argsort on an HPE EX system

Recent Journal Paper on using Chapel for calibrating hydrologic models

- Marjan Asgari et al, "Development of a knowledge-sharing parallel computing approach for calibrating distributed watershed hydrologic models", Environmental Modeling and Software.
- They report super-linear speedup







ARKOUDA ARGSORT PERFORMANCE

HPE Apollo (May 2021)



- HDR-100 Infiniband network (100 Gb/s)
- 576 compute nodes
- 72 TiB of 8-byte values
- ~480 GiB/s (~150 seconds)

HPE Cray EX (April 2023)

- Slingshot-11 network (200 Gb/s)
- 896 compute nodes
- 28 TiB of 8-byte values
- ~1200 GiB/s (~24 seconds)

HPE Cray EX (May 2023)

- Slingshot-11 network (200 Gb/s)
- 8192 compute nodes
- 256 TiB of 8-byte values
- ~8500 GiB/s (~31 seconds)

A notable performance achievement in ~100 lines of Chapel

Arkouda Argsort Performance



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CHAPEL EXECUTION MODEL AND TERMINOLOGY: LOCALES

- Locales can run tasks and store variables
 - Each locale executes on a "compute node" on a parallel system
 - User specifies number of locales on executable's command-line

| prompt> ./myC | hapelProgram | numLocales | =4 # or '-nl | 4' Fo | ur nodes/CPUs |
|----------------|--------------|----------------------|------------------------|----------|---------------|
| Locales array: | locale 0 | locale 1 | locale 2 | locale 3 | |
| | User's co | ode starts running a | is a single task on lo | cale O | |

hello-dist-node-names.chpl

const numTasks = here.numPUs();
coforall tid in 1..numTasks do
writef("Hello from task %n 1 %n on %s\n",
 tid, numTasks, here.name);
 what's my locale's name?



Hello from task 3 of 4 on n1032

Hello from task 2 of 4 on n1032

| <pre>> chpl hello-dist-node-names.chpl</pre> | | | | | | | | | | |
|---|------|------|---|----|---|----|-------|--|--|--|
| > ./hello-dist-node-names | | | | | | | | | | |
| Hello | from | task | 1 | of | 4 | on | n1032 | | | |
| Hello | from | task | 4 | of | 4 | on | n1032 | | | |
| Hello | from | task | 3 | of | 4 | on | n1032 | | | |
| Hello | from | task | 2 | of | 4 | on | n1032 | | | |

So far, this is a shared-memory program

Nothing refers to remote locales, explicitly or implicitly

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TASK-PARALLEL "HELLO WORLD" (DISTRIBUTED VERSION)





TASK-PARALLEL "HELLO WORLD" (DISTRIBUTED VERSION)



INTRODUCTION TO CHAPEL

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2D HEAT DIFFUSION EXAMPLE

See <u>https://go.lbl.gov/cuf23-repo</u> for more info and for example code.

See 'heat_2D.*.chpl' in the Chapel examples

- 'heat_2D.chpl' shared memory parallel version that runs in locale 0
- 'heat_2D_dist.chpl' parallel and distributed version that is the same as 'heat_2D.chpl' but with distributed arrays
- 'heat_2D_dist_buffers.chpl' parallel and distributed version that copies to neighbors landing pad and then into local halos

PARALLEL HEAT DIFFUSION IN HEAT_2D.CHPL



• 2D heat diffusion PDE



Simplified form for below assume $\Delta x = \Delta y$, and let $\alpha = \nu \Delta t / \Delta x^2$

• Solving for next temperatures at each time step using finite difference method

$$u_{i,j}^{n+1} = u_{i,j}^n + \alpha \left(u_{i+1,j}^n + u_{i-1,j}^n - 4u_{i,j}^n + u_{i,j+1}^n + u_{i,j-1}^n \right)$$

• All updates in a timestep can be done in parallel

• Output is the mean and standard deviation of all the values and time to solution

DISTRIBUTED AND PARALLEL HEAT DIFFUSION IN HEAT_2D_DIST.CHPL



• Declaring 'u' and 'un' arrays

```
const indices = {0..<nx, 0..<ny}
var u: [indices] real;</pre>
```

• Declaring 'u' and 'un' arrays as distributed (e.g., 2x2 distribution is shown)

```
const indices = {0..<nx, 0..<ny},
        INDICES = Block.createDomain(indices);
var u: [INDICES] real;
```

• Reads that cross the distribution boundary will result in a remote get

PARALLELISM SUPPORTED BY CHAPEL

Synchronous parallellism

- 'coforall', distributed memory parallelism across processes/locales with 'on' syntax
- 'coforall', shared-memory parallelism over threads
- 'cobegin', executes all statements in block in parallel

Asynchronous parallelism

- 'begin', creates an asynchronous task
- 'sync' and 'atomic' vars for task coordination
- spawning subprocesses

Higher-level parallelism abstractions

- 'forall', data parallelism and iterator abstraction
- 'foreach', SIMD parallelism
- 'scan', operations such as cumulative sums
- 'reduce', operations such as summation





LEARNING OBJECTIVES FOR TODAY'S CHAPEL TUTORIAL

- Compile and run Chapel programs
- Familiarity with the Chapel execution model including how to run codes in parallel on a single node, across nodes, and both
- Learn Chapel concepts by compiling and running provided code examples
 - Serial code using map/dictionary, (k-mer counting from bioinformatics)
 - Parallelism and locality in Chapel
 - Distributed parallelism and 1D arrays, (processing files in parallel)
 - Distributed parallelism and 2D arrays, (heat diffusion problem will see in UPC++ and CAF)
 - Distributed parallel image processing, (coral reef diversity example)
 - GPU parallelism (stream example)
- Where to get help and how you can participate in the Chapel community






Coarray Fortran Tutorial

Damian Rouson Computer Languages & System Software

Hosted by ECP, NERSC, and OLCF, 26-27 July 2023





- Introduction to Coarray Fortran ("CAF")
 - Why Fortran Matters
 - SPMD parallel execution
 - PGAS data structures & RMA
- Heat Conduction Solver
 - Compiling and running it
 - Understanding it

Why Fortran Matters



Intermediate Complexity Atmospheric Research (ICAR) Model Courtesy of Ethan Gutmann, NCAR

Weather & Climate



U.S. Nuclear Regulatory Commission File Photo

Nuclear Energy



FUN3D Mesh Adaptation for Mars Ascent Vehicle, Courtesy of Eric Nielsen & Ashley Korzun, NASA Langley

Aerospace

Why Fortran Matters



Intermediate Complexity Atmospheric Research (ICAR) Model Courtesy of Ethan Gutmann, NCAR

Weather & Climate



U.S. Nuclear Regulatory Commission File Photo

Nuclear Energy



FUN3D Mesh Adaptation for Mars Ascent Vehicle, Courtesy of Eric Nielsen & Ashley Korzun, NASA Langley

Aerospace

CAF Philosophy

"The underlying philosophy of our design is to make the smallest number of changes to the language required to obtain a robust and efficient parallel language without requiring the programmer to learn very many new rules."

Reid, J., & Numrich, R. W. (2007). Co-arrays in the next Fortran standard. *Scientific Programming*, *15*(1), 9-26.

Seminal paper:

Numrich, R. W., & Reid, J. (1998, August). Co-Array Fortran for parallel programming. In *ACM SIGPLAN Fortran Forum* (Vol. 17, No. 2, pp. 1-31). New York, NY, USA: ACM.



Single Program Multiple Data



cd fortran make run-hi

Single Program Multiple Data (SPMD) parallel execution

- Synchronized launch of multiple "images" (process/threads/ranks)
- Asynchronous execution except where program explicitly synchronizes
- Error termination or synchronized normal termination



Compiling and Running hi.f90



Compiling and Running hi.f90





7

Image 1

- 1 program main
- 2 implicit none
- 3 print *,"Hello from image ", this_image(), "of", num_images()
 4 end program

💿 rouson — vim hi.f90 — 67×5

- 1. After the creation of a fixed number of images, each image's first "segment" (sequence of statements) executes.
- Image control statements totally order segments executed by a single image and partially order segments executed by separate images.



7

Image 1

• • •

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💼 rouson — vim hi.f90 — 67×5



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|---|---|---|---|--|--|
| • | • | • | | | |
| | | | _ | | |

🛅 rouson — vim hi.f90 — 67×5

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Partitioned Global Address Space (PGAS)

Coarrays:

- Distributed data structures greeting
- Facilitate Remote Memory Access (RMA) line 15



cd fortran make run-hello

```
cuf23-tutorial — vim hello.f90 — 74×21
. . .
 1 program main
 2
      !! One-sided communication of distributed greetings
 3
     implicit none
     integer, parameter :: max_greeting_length=64, writer = 1
 4
     integer image
 5
 6
      character(len=max_greeting_length) :: greeting[*] ! scalar coarray
 7
 8
      associate(me => this_image(), ni=>num_images())
 9
10
        write(greeting,*) "Hello from image",me,"of",ni ! local (no "[]")
11
        sync all ! image control
12
13
        if (me == writer) then
14
          do image = 1, ni
            print *, greeting[image] ! one-sided communication: "get"
15
16
          end do
17
        end if
18
      end associate
19
20 end program
```

Compiling & Running hello.f90





Compiling & Running hello.f90





Compiling and Running the Heat Equation Solver





Compiling and Running the Heat Equation Solver





Heat Equation



cd fortran make run-heat-equation

$$\frac{\partial T}{\partial t} = \alpha \nabla^2 T$$

$$\{T\}^{n+1} = \{T\}^n + \Delta t \cdot \alpha \cdot \nabla^2 \{T\}^n$$

T = T + dt * alpha * .laplacian. T

Heat Equation

 \sim



cd fortran make run-heat-equation

$$\frac{\partial T}{\partial t} = \alpha \nabla^2 T$$

$$\{T\}^{n+1} = \{T\}^n + \Delta t \cdot \alpha \cdot \nabla^2 \{T\}^n$$

$$[\mathbf{T}] = [\mathbf{T}] + d\mathbf{t} * alpha * .laplacian.$$



Heat Equation



cd fortran make run-heat-equation

$$\frac{\partial T}{\partial t} = \alpha \nabla^2 T$$

$$\{T\}^{n+1} = \{T\}^n + \Delta t \cdot \alpha \cdot \nabla^2 \{T\}^n$$

$$\mathbf{T} = \mathbf{T} + dt \mathbf{*} \text{ alpha} \mathbf{*} \text{ laplacian. } \mathbf{T}$$

$$\mathbf{I}$$

$$\mathbf{I}$$

pure user-defined operators

Class Diagram





Halo Exchange





```
116 real(rkind), allocatable :: halo_x(:,:)[:]
117 integer, parameter :: west=1, east=2
```

```
134 me = this_image()
135 num_subdomains = num_images()
137 my_nx = nx/num_subdomains + merge(1, 0, me <= mod(nx, num_subdomains))
232 subroutine exchange_halo(self)
233 class(subdomain_2D_t), intent(in) :: self
234 if (me>1) halo_x(east,:)[me-1] = self%s_(1,:)
235 if (me<num_subdomains) halo_x(west,:)[me+1] = self%s_(my_nx,:)
236 end subroutine</pre>
```

Loop-Level Parallelism



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| TAU: ParaProf: Statistics for: node 0 - /home/tutorial/SRC/demo/matcha | | | - | • • |
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| Name | Exclu | Inclu 🔻 | Calls | Chil |
| TAU application | 0 | 1.516 | 1 | 1 |
| taupreload_main | 0.801 | 1.516 | 10 | 51,499 |
| è ■[CONTEXT] taupreload_main | 0 | 0.811 | 27 | (|
| [SUMMARY]subdomain_2d_m_MOD_laplacian [{/home/tutorial/SRC/demo/matcha/example/heat-equation.f90}] | 0.6 | 0.6 | 20 | (|
| [SAMPLE]subdomain_2d_m_MOD_laplacian [{/home/tutorial/SRC/demo/matcha/example/heat-equation.f90} {188} |] 0.54 | 0.54 | 18 | |
| [SAMPLE]subdomain_2d_m_MOD_laplacian [{/home/tutorial/SRC/demo/matcha/example/heat-equation.f90} {183} |] 0.03 | 0.03 | 1 | (|
| [SAMPLE]subdomain_2d_m_MOD_laplacian [{/home/tutorial/SRC/demo/matcha/example/heat-equation.f90} {187} |] 0.03 | 0.03 | 1 | (|
| [SAMPLE]subdomain_2d_m_MOD_copy [{/home/tutorial/SRC/demo/matcha/example/heat-equation.f90} {217}] | 0.06 | 0.06 | 2 | (|
| [SAMPLE]subdomain_2d_m_MOD_add [{/home/tutorial/SRC/demo/matcha/example/heat-equation.f90 } {212 }] | 0.06 | 0.06 | 2 | (|
| [SAMPLE]subdomain_2d_m_MOD_multiply [{/home/tutorial/SRC/demo/matcha/example/heat-equation.f90} {207}] | 0.03 | 0.03 | 1 | |
| [SAMPLE] raw_write [{unix.c} {0}] | 0.03 | 0.03 | 1 | |
| □ [SAMPLE]tls_get_addr [{/usr/lib64/ld-2.26.so} {0}] | 0.03 | 0.03 | 1 | |
| MPI_Win_lock() | 0.363 | 0.3632 | 20,481 | |
| MPI_Barrier() | 0.21 | 0.21 | 12 | (|
| MPI_Finalize() | 0.094 | 0.094 | 1 | 1 |
| | 0.018 | 0.0182 | 20,481 | (|
| | 0.017 | 0.0172 | 20,480 | (|
| | 0.01 | 0.01 | 1 | (|
| MPI Collective Sync | 0.002 | 0.002 | 2 | (|
| MPI_Comm_aup() | 0 | 0.001 | 1 | 1 |
| MPI_win_create() | 0 | 0 | 1 | (|

Comments



- Coarray Fortran began as a syntactically small extension to Fortran 95:
 - Square-bracketed "cosubscripts" distribute & communicate data
- Integration with other features:
 - -Array programming: colon subscripts
 - -OOP: distributed objects
 - Minimally invasive:
 - Drop brackets when not communicating
- Communication is explicit:
- -Use brackets when communicating

| | • |
|--|---|
| | |
| • • • | Desktop — vim pgas.f90 — 56×15 |
| orogram main | |
| implicit n | ione |
| type foo | |
| integer | :: bar=2 |
| end type | |
| integer, p | arameter :: local_size=5 |
| type(foo) | <pre>:: object(local_size)[*]=foo()</pre> |
| associate(| <pre></pre> |
| if (n<3) | error stop "Insufficient number of images." |
| sync all | |
| if (me <n< td=""><td>) $object(1:2) = object(3:4) [me+1]$</td></n<> |) $object(1:2) = object(3:4) [me+1]$ |
| if (me== | <pre>1) object(5)[2] = object(5)[3]</pre> |
| end associ | .ate |
| end program | |
| | |





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UPC++: An Asynchronous RMA/RPC Library for Distributed C++ Applications

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GASNet-EX









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What does UPC++ offer?

Asynchronous behavior

- RMA:
 - Get/put to a remote location in another address space
 - Low overhead, zero-copy, one-sided communication.
- RPC: Remote Procedure Call:
 - Moves computation to the data

Design principles for performance

- All communication is syntactically explicit
- All communication is asynchronous: futures and promises
- Scalable data structures that avoid unnecessary replication



Some motivating applications

Many applications involve asynchronous updates to irregular data structures

- Adaptive meshes
- Sparse matrices
- Hash tables and histograms
- Graph analytics
- Dynamic work queues

Irregular and unpredictable data movement:

- Space: Pattern across processors
- Time: When data moves
- Volume: Size of data







Graph analytics



Some motivating system trends

The first exascale systems appeared in 2022

- Cores per node is growing
- Accelerators (e.g. GPUs) are becoming more important
- Latency is not improving

Need to reduce communication costs in software

- Overlap communication to hide latency
- Reduce memory using smaller, more frequent messages
- Minimize software overhead
- Use simple messaging protocols (RDMA)









Reducing communication overhead

Let each process directly access another's memory via a global pointer Communication is **one-sided** – there is no "receive" operation

- No need to match sends to receives
- No unexpected messages
- No need to guarantee message ordering



- All metadata provided by the initiator, rather than split between sender and receiver
- Supported in hardware through RDMA (Remote Direct Memory Access)

Looks like shared memory: shared data structures with asynchronous access


One-sided GASNet-EX vs one- and two-sided MPI

Four distinct network hardware types

The performance of one-sided GASNet-EX matches or exceeds that of MPI RMA and message-passing:

- 8-byte Put latency 19 52% better
- 8-byte Get latency 16 49% better
- Better flood bandwidth efficiency: often reaching same or better peak at ½ or ¼ the transfer size

8-Byte RMA Operation Latency (one-at-a-time)



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Perlmutter Phase-I results collected July 2022, all others collected April 2023. GASNet-EX tests were run using then-current GASNet library and its tests. MPI tests were run using then-current center default MPI version and Intel MPI Benchmarks. All tests use two nodes and one process per node. For details see LCPC'18 <u>doi.org/10.25344/S4QP4W</u> and PAW-ATM'22 <u>doi.org/10.25344/S40C7D</u> See also: <u>gasnet.lbl.gov/performance</u>

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A Partitioned Global Address Space programming model

Global Address Space

- Processes may read and write *shared segments* of memory
- Global address space = union of all the shared segments

Partitioned

- Global pointers to objects in shared memory have an affinity to a particular process
- Explicitly managed by the programmer to optimize for locality
- In conventional shared memory, pointers do not encode affinity





The PGAS model

Partitioned Global Address Space

- Support global memory, leveraging the network's RDMA capability
- Distinguish private and shared memory
- Separate synchronization from data movement

Languages that provide PGAS: **Chapel**, **Co-Array Fortran (Fortran 2008)**, UPC, Titanium, X10

Libraries that provide PGAS: OpenSHMEM, Co-Array C++, Global Arrays, DASH, MPI-RMA

This presentation is about UPC++, a C++ library developed at Lawrence Berkeley National Laboratory



Execution model: SPMD

Like MPI and Coarray Fortran, UPC++ uses a SPMD model of execution, where a fixed number of processes run the same program

```
int main() {
upcxx::init();
cout << "Hello from " << upcxx::rank_me() << endl;
upcxx::barrier();
if (upcxx::rank_me() == 0) cout << "Done." << endl;
upcxx::finalize();</pre>
```





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}

Global pointers

Global pointers are used to create logically shared but physically distributed data structures

Parameterized by the type of object it points to, as with a C++ (raw) pointer: e.g. <u>global_ptr</u><double>, <u>global_ptr</u><Node>



Global vs raw pointers and affinity

The affinity identifies the process that created the object

Global pointer carries both an address and the affinity for the data

Raw C++ pointers (e.g. Node*) can be used on a process to refer to objects in the global address space that have affinity to that process



How does UPC++ deliver the PGAS model?

UPC++ uses a "compiler-free," library approach

• UPC++ leverages C++ standards, needs only a standard C++ compiler



Relies on GASNet-EX for low-overhead communication

- Efficiently utilizes network hardware, including RDMA
- Provides Active Messages on which UPC++ RPCs are built
- Enables portability (laptops to supercomputers)

Designed for interoperability

- Same process model as MPI, enabling hybrid applications
- On-node compute models (e.g. OpenMP, CUDA, HIP, Kokkos) can be mixed with UPC++ as in MPI+X



UPC++ on top of GASNet



Cray XC40 system

Two processor partitions:

- Intel Haswell (2 x 16 cores per node)
- Intel KNL (1 x 68 cores per node)



Round-trip Put Latency (lower is better) Flood Put Bandwidth (higher is better) Data collected on Cori Haswell (https://doi.org/10.25344/S4V88H)



Asynchronous communication (RMA)

By default, all communication operations are split-phased

- Initiate operation
- Wait for completion

A future holds a value and a state: ready/not-ready



Remote procedure call (RPC)

Execute a function on another process, sending arguments and returning an optional result

- 1. Initiator injects the RPC to the target process
- 2. Target process executes fn(arg1, arg2) at some later time determined at the target
- 3.Result becomes available to the initiator via the future

Many RPCs can be active simultaneously, hiding latency



Hands-on: 2D heat diffusion

| $u_{i,j}^{n+1} = u_{i,j}^n + \alpha \left(u_{i+1,j}^n + u_{i-1,j}^n - 4u_{i,j}^n + u_{i,j+1}^n + u_{i,j}^n \right)$ | -1) |
|--|-----|
|--|-----|

Everything needed for the hands-on activities is at: <u>https://go.lbl.gov/CUF23</u>

Online materials include:

- Module info for NERSC Perlmutter, OLCF Frontier, and other machines
- Download links to install UPC++

Once you have set up your environment, copied the tutorial materials, and changed to the cuf23/upcxx directory:

